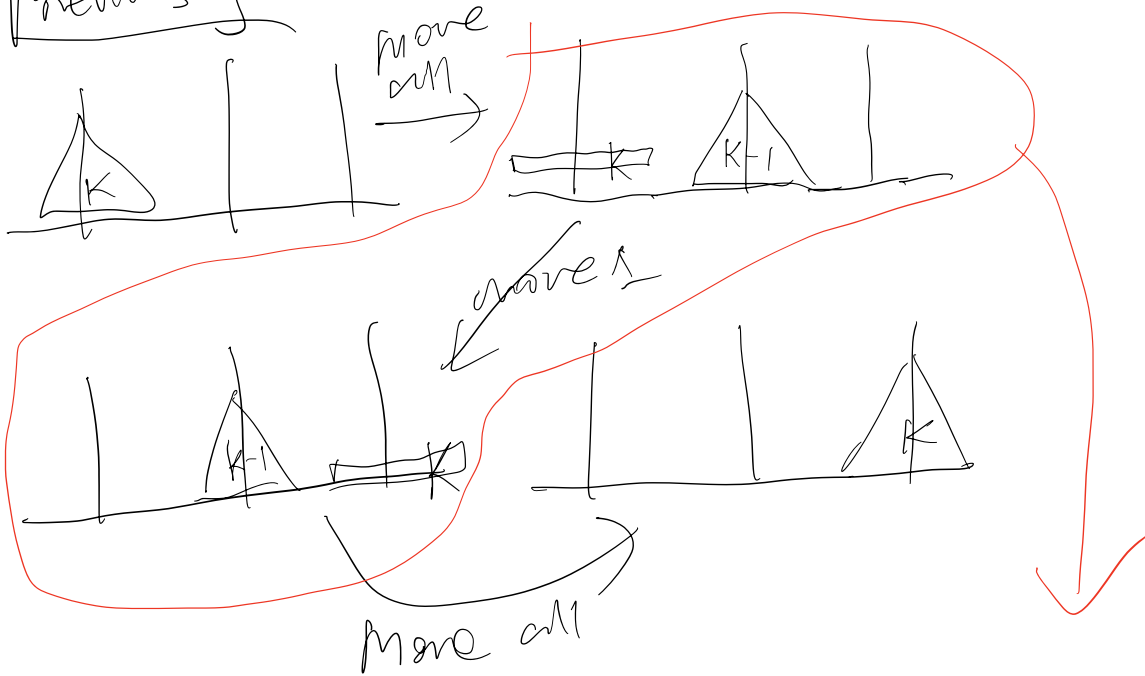


23.1

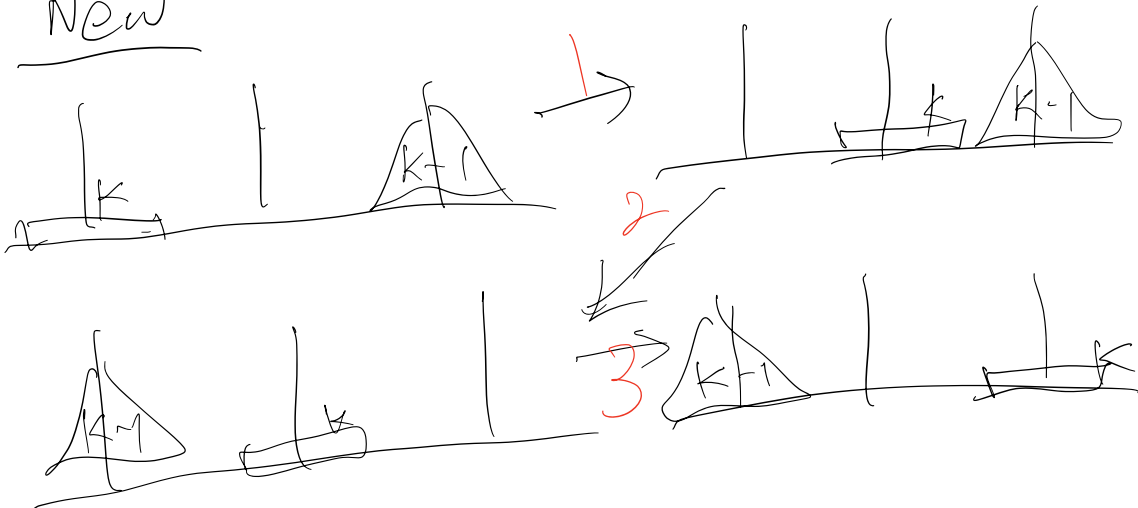
```
20 void solve(long k, char source, char dest, char placeholder) {  
21     if (k == 1) {  
22         move(k, source, dest);  
23         return;  
24     }  
25  
26     solve(k - 1, source, placeholder, dest);  
27     move(k, source, dest);  
28     solve(k - 1, placeholder, dest, source);  
29 }
```

Needs to be repeated.

Previously



New

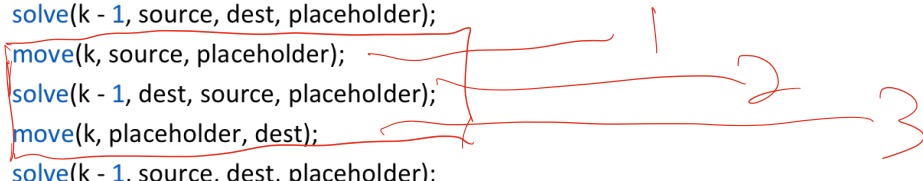


```

void solve(long k, char source, char dest, char placeholder) {
    if (k == 1) {
        move(k, source, placeholder);
        move(k, placeholder, dest);
        return;
    }
    solve(k - 1, source, dest, placeholder);
    move(k, source, placeholder);
    solve(k - 1, dest, source, placeholder);
    move(k, placeholder, dest);
    solve(k - 1, source, dest, placeholder);
}

```

A → B  
B → C



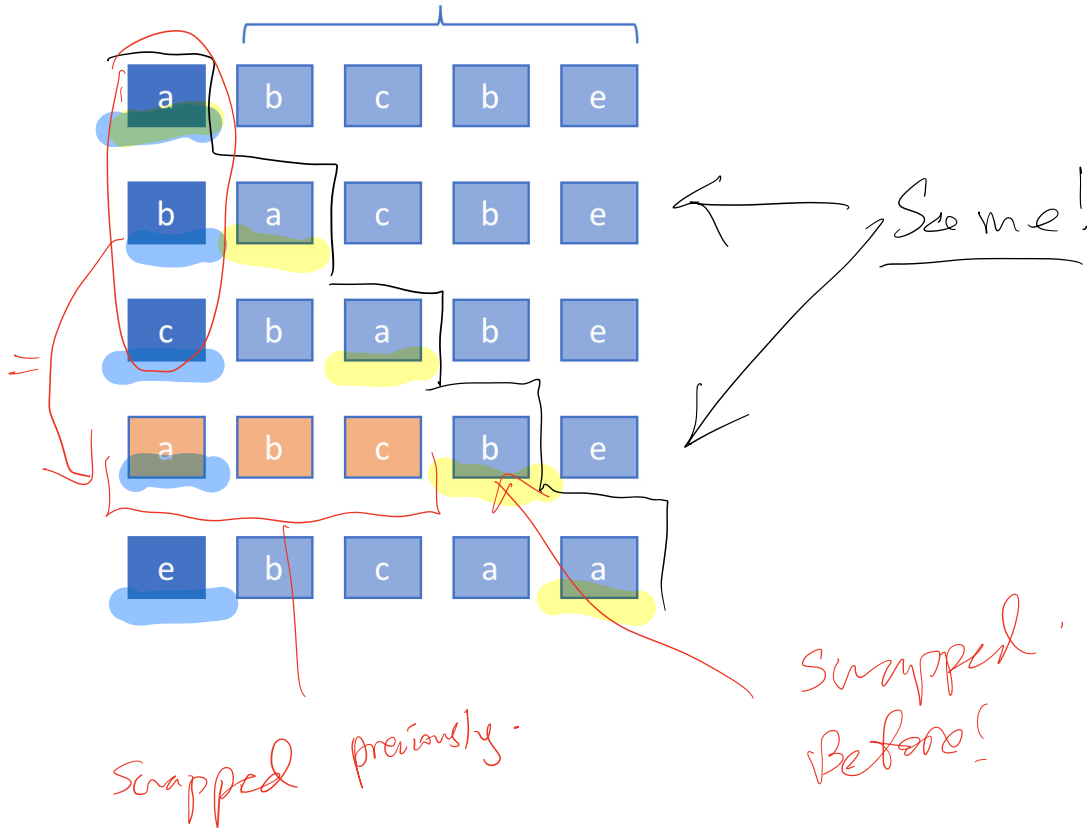
$$T(k) = \begin{cases} 3T(k-1) + 2, & \text{if } k \geq 1 \\ 2, & k = 1 \end{cases}$$

So, analysis follows ...

$$O(3^k)$$

24.1

### permute recursively



```
1  /**
2   * Fix a[0]..a[curr - 1] but permute characters a[curr]..a[len - 1]
3   * Print out each permutation.
4   *
5   * @param[in,out] a    The array to permute
6   * @param[in]      len The size of the array
7   * @param[in]      curr The starting index at which we will permute
8   *
9   * @post The string a remains unchanged
10  */
11 void permute(char a[], size_t len, size_t curr) {
12     if (curr == len - 1) {
13         cs1010_println_string(a);
14         return;
15     }
16
17     for (size_t i = curr; i < len; i += 1) {
18         swap(a, curr, i);
19         permute(a, len, curr + 1);
20         swap(a, i, curr);
21     }
22 }
```

25.1

```

30  /**
31  * Checks if any queen from row 0 to last_row (inclusive)
32  * that clashes with each other, diagonally.
33  *
34  * @param[in] queens  The array containing the representation
35  *                    of the queens.
36  * @param[in] last_row The last row until which we check for
37  *                    clashes.
38  *
39  * @pre  0 <= last_row <= n-1
40  * @return true if there are two queens that clash with each other.
41  */
42  bool threaten_each_other_diagonally(char queens[], size_t last_row) {
43      for (size_t begin_row = 0; begin_row <= last_row; begin_row += 1) {
44          if (has_a_queen_in_diagonal(queens, begin_row, last_row)) {
45              return true;
46          }
47      }
48      return false;
49  }

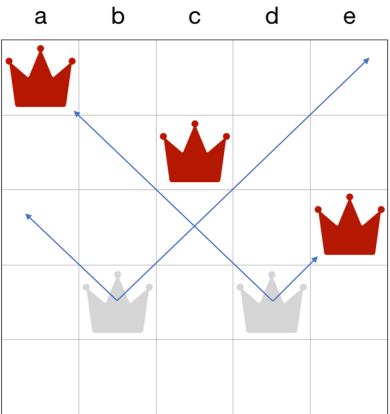
```

**N-Queens Solution: Version 3**

```

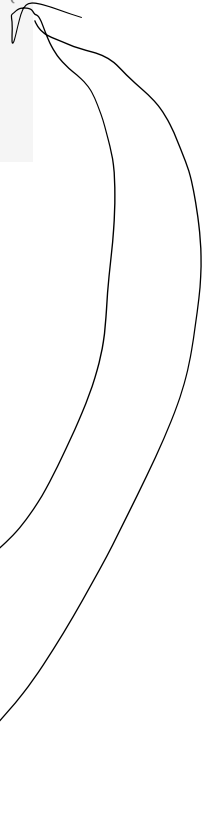
1  /**
2  * Search for all possible queens placement from row to n-1,
3  * given the queens placement from row 0 to row-1.
4  *
5  * @param[in] queens  The string representation of queens
6  *                    placement.
7  * @param[in] n        The size of the chess board
8  * @param[in] row      The last row where the queens positions
9  *                    have been fixed.
10  */
11  bool nqueens(char queens[], long n, long row) {
12      if (row == n - 1) {
13          if (!threaten_each_other_diagonally(queens, n - 1)) {
14              cs1010_println_string(queens);
15              return true;
16          }
17          return false;
18      }
19
20      for (long i = row; i < n; i++) {
21          swap(queens, row, i);
22          if (!threaten_each_other_diagonally(queens, row)) {
23              if (nqueens(queens, n, row + 1)) {
24                  return true;
25              }
26          }
27          swap(queens, row, i);
28      }
29      return false;
30  }

```



Already "safe."  
No need to check again.

Only check the new row against the queens above.



25.2

```
void permute(char a[], size_t len, size_t curr) {  
  if (curr == len-1) {  
    if (a[curr] != a[curr-1]) {  
      cs1010_println_string(a);  
    }  
    return;  
  }  
  for (size_t i = curr; i < len; i += 1) {  
    if (!appear_before(a, curr, i) && a[i] != a[curr-1]) {  
      swap(a, curr, i);  
      permute(a, len, curr + 1);  
      swap(a, i, curr);  
    }  
  }  
}
```

↙  
Before swap  
a[curr]