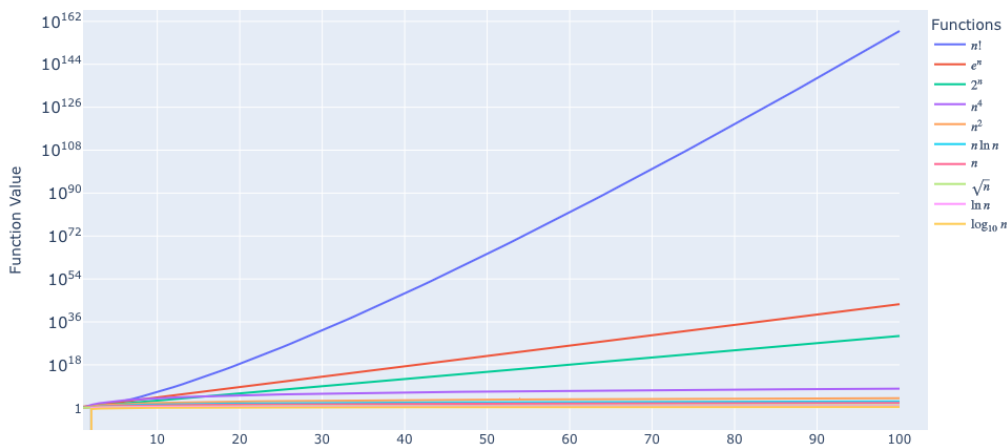


$$2 < e \Rightarrow 2^n < e^n$$

$$\left. \begin{array}{l} n! = 1 \times 2 \times \dots \times n \\ e^n = e \times e \times \dots \times e \end{array} \right\} \downarrow < 2^n < e^n < n!$$

<https://eric-han.com/teaching/AY2425S1/CS1010/P20.1.html>

Comparison of Various Mathematical Functions



Why do we care about complexity?

What is the Big-O running time of the following code, in terms of n ?

a) $O(n \times n/2) \sim O(n^2)$

```

1 for (long i = 0; i < n; i += 1) {
2   for (long j = 0; j < n; j += 2) {
3     cs1010_println_long(i + j);
4   }
5 }

```

Handwritten annotations: Above the first loop, "0, 1, 2... n-1" and "n" are written. Above the second loop, "0, 2, ..." and " $\lceil n/2 \rceil$ " are written.

b) $O(\log(n) \times \log(n)) \sim O(\log^2 n)$

```

1 for (long i = 1; i < n; i *= 2) {
2   for (long j = 1; j < n; j *= 2) {
3     cs1010_println_long(i + j);
4   }
5 }

```

Handwritten annotations: Above the first loop, "1, 2, 4, ..." and " $\lceil \log_2(n) \rceil$ " are written. Above the second loop, "1, 2, 4, ..." and " $\lceil \log_2(n) \rceil$ " are written.

c) $O(1 + 2 + \dots + n) = O\left(\sum_{i=1}^n i\right) = \frac{n(n+1)}{2} = O(n^2)$ *Arith Series*

```

1 for (long j = 0; j < n; j += 1) {
2   for (long i = 0; i < j; i += 1) {
3     cs1010_println_long(i + j);
4   }
5 }

```

Handwritten annotations: A diagram shows nested loops. The outer loop has values "0, 1, 2... n-1" and "n". The inner loop has values "0, 1, 2... n-1", "0, 1, 2", "0, 1", and "0".

d) $O(2^1 + 2^2 + \dots + 2^n) = O\left(\sum_{i=1}^n 2^i\right) = 2(2^n - 1) = O(2^n)$ *Geo. Series*

```

1 long k = 1;
2 for (long j = 0; j < n; j += 1) {
3   k *= 2;
4   for (long i = 0; i < k; i += 1) {
5     cs1010_println_long(i + j);
6   }
7 }

```

Handwritten annotations: A diagram shows nested loops. The outer loop has values "0, 1, 2... n-1" and "n". The inner loop has values "0, 1, 2, 3" (labeled "k=4", "4=2^2") and "0, 1" (labeled "k=2", "2=2^1").

a) Express the running time of the following function as a recurrence relation: *Recursion*

```

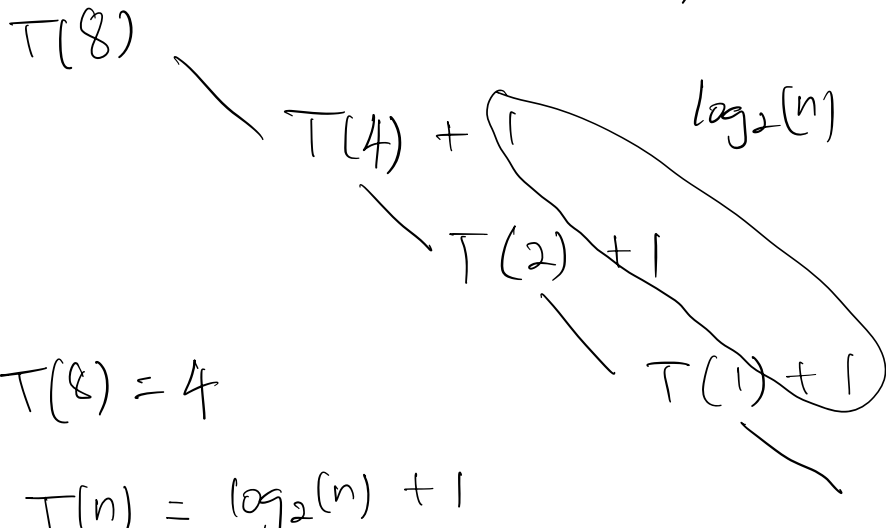
1 void foo(long n) {
2   if (n == 1) {
3     return 1;
4   }
5   return foo(n/2) + 2;
6 }

```

$$foo(n) = \begin{cases} 1 & \text{if } n = 1 \\ foo(n/2) + 2 & \text{else} \end{cases}$$

What is its running time?

$$T(n) = T(n/2) + 1, T(1) = 1$$



$$T(8) = 4$$

$$T(n) = \log_2(n) + 1$$

$$O(\log n)$$

b) Express the running time of the following function as a recurrence relation:

```

1 void foo(long n) {
2   if (n == 1) {
3     return 1;
4   }
5   for (long i = 0; i < n; i += 1) {
6     cs1010_println_long(i);
7   }
8   return foo(n - 1);
9 }

```

$0, 1, \dots, n-1$ n

$T(n) = T(n-1) + n$

What is its running time?

$T(1) = 1$

